

# A Modified Verification Method for Linear Systems

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This talk is concerned with the problem of verifying the accuracy of approximate solutions of linear systems. Let  $\mathbb{R}$  be the set of real numbers. For  $A \in \mathbb{R}^{n \times n}$ , the comparison matrix of  $A$  is denoted by  $\langle A \rangle$ .

For a linear system

$$Ax = b, \quad A \in \mathbb{R}^{n \times n}, \quad b \in \mathbb{R}^n, \quad (1)$$

we can efficiently obtain a computed solution by some numerical algorithm. In general, however, we do not know how accurate the computed solution is. For this purpose, several verification methods have been proposed (cf. e.g. [1,2]).

In this talk we propose an efficient method of calculating a componentwise error bound of the computed solution of (1), which is based on the following Rump's theorem:

**Theorem 1** (Rump [1, Theorem 2.1]). *Let  $A \in \mathbb{R}^{n \times n}$  and  $b, \tilde{x} \in \mathbb{R}^n$  be given. Assume  $v \in \mathbb{R}^n$  with  $v > \mathbf{0}$  satisfies  $u := \langle A \rangle v > \mathbf{0}$ . Let  $\langle A \rangle = D - E$  denote the splitting of  $\langle A \rangle$  into the diagonal part  $D$  and the off-diagonal part  $-E$ , and define  $w \in \mathbb{R}^n$  by*

$$w_k := \max_{1 \leq i \leq n} \frac{G_{ik}}{u_i} \quad \text{for } 1 \leq k \leq n,$$

where  $G := I - \langle A \rangle D^{-1} = ED^{-1} \geq O$ . Then  $A$  is nonsingular, and

$$|A^{-1}b - \tilde{x}| \leq (D^{-1} + vw^T)|b - A\tilde{x}|. \quad (2)$$

In particular, the method based on Theorem 1 is implemented in the routine `verifylss` in INTLAB Version 7 [3].

We modify Theorem 1 as follows:

**Theorem 2.** Let  $A, b, \tilde{x}, u, v, w$  be defined as in Theorem 1. Define  $D_s := \text{diag}(s)$  where  $s \in \mathbb{R}^n$  with

$$s_k := u_k w_k \quad \text{for } 1 \leq k \leq n.$$

Then

$$|A^{-1}b - \tilde{x}| \leq (D^{-1} + vw^T)(I + D_s)^{-1}|b - A\tilde{x}|. \quad (3)$$

*Proof.* From the definition of  $u$  and  $w$ , it holds

$$I - \langle A \rangle D^{-1} \leq uw^T.$$

Since  $\text{diag}(I - \langle A \rangle D^{-1}) = \mathbf{0}$ , we have

$$I - \langle A \rangle D^{-1} + D_s \leq uw^T$$

and

$$I + D_s \leq \langle A \rangle D^{-1} + uw^T. \quad (4)$$

From the assumption,  $A$  is an  $H$ -matrix, so that  $\langle A \rangle^{-1} \geq O$ . Multiplying (4) from the left by  $\langle A \rangle^{-1}$  yields

$$\langle A \rangle^{-1}(I + D_s) \leq D^{-1} + \langle A \rangle^{-1}uw^T = D^{-1} + vw^T.$$

Since  $(I + D_s)^{-1} \geq O$ , we have

$$\langle A \rangle^{-1} \leq (D^{-1} + vw^T)(I + D_s)^{-1}. \quad (5)$$

Using  $|(A)^{-1}| \leq \langle A \rangle^{-1}$  and (5),

$$\begin{aligned} |A^{-1}b - \tilde{x}| &\leq |A^{-1}||b - A\tilde{x}| \leq \langle A \rangle^{-1}|b - A\tilde{x}| \\ &\leq (D^{-1} + vw^T)(I + D_s)^{-1}|b - A\tilde{x}|, \end{aligned}$$

which proves the theorem.  $\square$

Theorem 2 always gives better bounds than Theorem 1 since  $(I + D_s)^{-1} \leq I$ . Numerical results will be shown to illustrate the efficiency of the proposed theorem. A generalization of Theorem 2 will be presented.

### References:

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