# On the maximum relative error. when computing $x^{n}$ in floating-point arithmetic 

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We improve the usual relative error bound for the computation of $x^{n}$ through iterated multiplications by $x$ in binary, precision- $p$, floating-point arithmetic. More precisely, we analyze the following algorithm.

$$
\begin{aligned}
& y \leftarrow x \\
& \text { for } k=2 \text { to } n \text { do } \\
& y \leftarrow \operatorname{RN}(x \cdot y) \\
& \text { end for } \\
& \text { return }(y)
\end{aligned}
$$

where RN means "round to nearest" (that is, $\operatorname{RN}(x \cdot y)$ is the result of the floating point multiplication $\mathrm{x} * \mathrm{y}$ in round-to-nearest mode).

We show the following result:
Theorem 1. Assume $p \geq 5$ (which holds in all practical cases). If

$$
n \leq \sqrt{2^{1 / 3}-1} \cdot 2^{p / 2}
$$

then

$$
\left|\widehat{x}_{n}-x^{n}\right| \leq(n-1) \cdot u \cdot x^{n}
$$

The obtained relative error bound $(n-1) \cdot u$ is only slightly better than the usual one $\gamma_{n-1}$, but it is simpler. We also discuss the more general problem of computing the product of $n$ terms.

## References:

[1] Muller et al., Handbook of Floating-Point Arithmetic, Birkhauser, 2010.
[2] Graillat, Lefevre and Muller, On the maximum relative error when computing $x^{n}$ in floating-point arithmetic, Technical report, 2013. Available at http://hal-ens-lyon.archives-ouvertes.fr/ensl-00945033

